
CMPSCI 311: Introduction to Algorithms

First Midterm Exam: Practice Exam

Name: _____ ID: _____

Instructions:

- Answer the questions directly on the exam pages.
- Show all your work for each question. Providing more detail including comments and explanations can help with assignment of partial credit.
- If the answer to a question is a number, *unless the problem says otherwise*, you may give your answer using arithmetic operations, such as addition, multiplication, “choose” notation and factorials (e.g., “ $9 \times 35! + 2$ ” or “ $0.5 \times 0.3 / (0.2 \times 0.5 + 0.9 \times 0.1)$ ” is fine).
- If you need extra space, use the back of a page.
- No books, notes, calculators or other electronic devices are allowed. Any cheating will result in a grade of 0.
- If you have questions during the exam, raise your hand.

Question	Value	Points Earned
1	10	
2	10	
3	10	
4	10	
5	10	
Total	50	

Question 1. (10 points) Indicate whether each of the following statements is TRUE or FALSE. No justification required.

1.1 (2 points): $\sum_{i=1}^n 1/i^2 = \Theta(n^4)$.

1.2 (2 points): A graph with n vertices and $n - 1$ edges is either disconnected or a tree.

1.3 (2 points): For every n there exists a directed graph on n vertices with $\Omega(n^2)$ edges that has a topological ordering.

1.4 (2 points): In a connected weighted graph, the edge with maximum weight is never in the minimum spanning tree.

1.5 (2 points): The recurrence $T(n) = 4T(n/2) + O(n)$ solves to $T(n) = \Theta(n^3)$.

Question 2. (10 points)

2.1 (5 points): Recall the scheduling problem where we have several task with lengths $t(i)$ and deadlines $d(i)$ and we want to order the tasks to minimize lateness where, if task i is completed at time $f(i)$, then lateness is defined as $L = \max_i \max(0, f(i) - d(i))$. Prove that ordering the intervals by their slack time, i.e., $d(i) - t(i)$ fails to find an optimal solution.

2.2 (5 points): On a stable matching instance, prove that if we run the Gale-Shapley algorithm twice, once with schools proposing and once with students proposing and we obtain the same matching, then the instance has a unique stable solution.

Question 3. (*10 points*) Alice is planning her course schedule for her time at UMass. There are n courses she must take and each course c_i can have pre-requisites P_i , which is a possibly empty set of courses. However, the department allows students to take a course and its prerequisites in the same semester. In other words, a course c_i can be taken in semester t if for all $c_j \in P_i$ the semester in which Alice takes c_j is at most t .

On the other hand, Alice can take at most 3 courses in a semester.

1. Prove that if the pre-requisite graph has a cycle of length 4, then there is no way for Alice to find a schedule satisfying all the pre-requisites.

2. Prove that if every course is involved in at most one cycle of length at most 3, then a valid schedule must exist.

Question 4. (10 points) Given two lists, L_1 of length n and L_2 of length m . We say that L_2 is a *subsequence* of L_1 if we can remove elements from L_1 to produce L_2 . This means that there exists indices $i_1 < \dots < i_m$ such that $L_1[i_j] = L_2[j]$ for each j . Design an algorithm that detects if L_2 is a subsequence of L_1 and outputs the indices i_1, \dots, i_m if L_2 is a subsequence of L_1 .

Question 5. (10 points) Suppose we have a complete k -ary tree with n leaves (suppose $n = k^d$ for some integer d). Each leaf v is associated with a weight $w(v)$. The weight of an internal node is defined to be the sum of the weights of all leaves that are descendants of this node. So the weight of the root r is $w(r) = \sum_{\text{leaves } v} w(v)$. Design and analyze an algorithm to compute the weight of every internal node.