Flux A Language for Programming High-Performance Servers

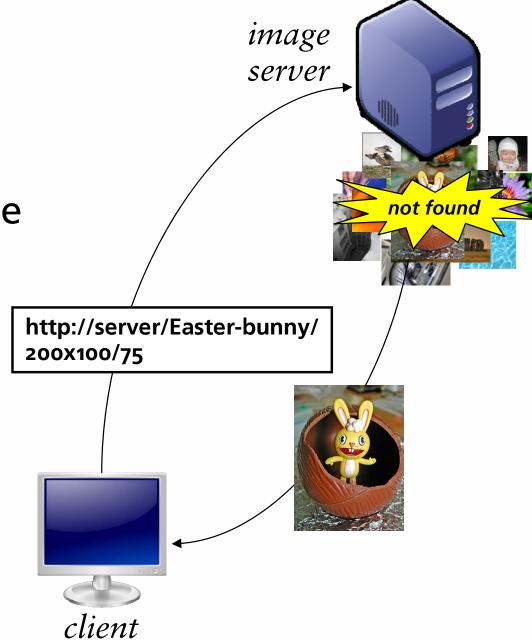
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Motivation: Image Server

Client

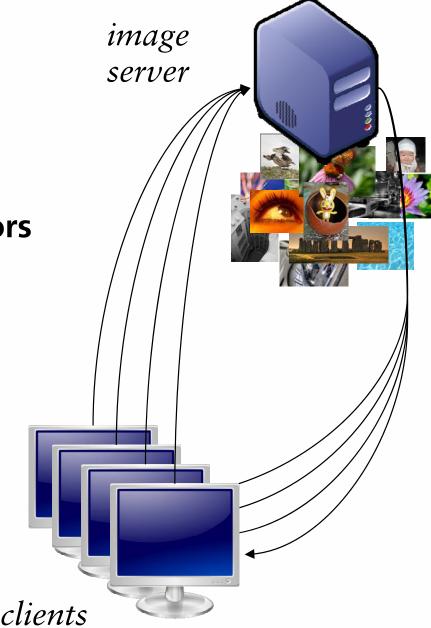
- Requests image @ desired quality, size
- Server
 - Images: RAW
 - Compresses to JPG
 - Caches requests
 - Sends to client





Problem: Concurrency

- Sequential fine until:
 - More clients
 - Bigger server
 - Multicores, multiprocessors
- One approach: threads
 - Limit reuse, risk deadlock, burden programmer
 - Complicate debugging
 - Mixes program logic & concurrency control





The Flux Programming Language

High-performance & deadlock-free concurrent programming w/ sequential components

Flux = Components + Flow + Atomicity

- Components = off-the-shelf C, C++, or Java
- Flow = path through components

Implicitly parallel

Atomicity = lightweight constraints

Compiler generates:

- Deadlock-free server
 - Runtime independent (threads, events, ...)
- Discrete event simulator



Outline

Intro to Flux: building a server

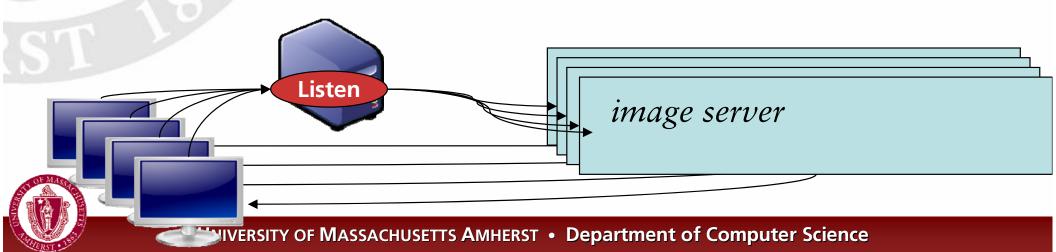
- Components
- Flows
- Atomicity
- Performance results
 - Server performance
 - Performance prediction
- Future work



Flux Server: Main

Source: one flow per connection
 Conceptually: separate thread
 Executes inside implicit infinite loop

source Listen \rightarrow Image;

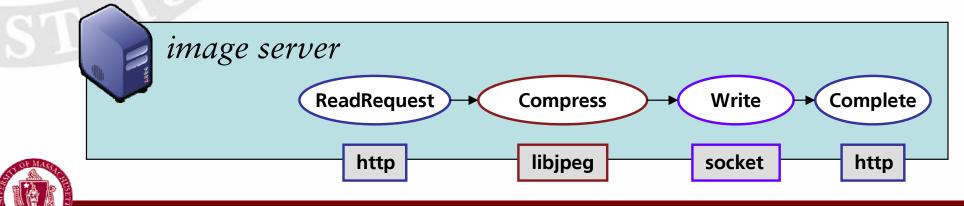


Flux Image Server

Basic image server requires:

- HTTP parsing (http)
- Socket handling (socket)
- JPEG compression (libjpeg)
- Single flow implements basic server:

Image = ReadRequest \rightarrow Compress \rightarrow Write \rightarrow Complete;



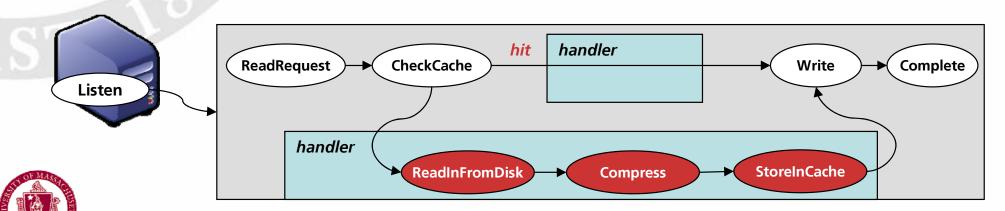
Adding Caching

Cache frequently requested images

Increase performance

- Direct data flow with predicates
 - Type test applied to output



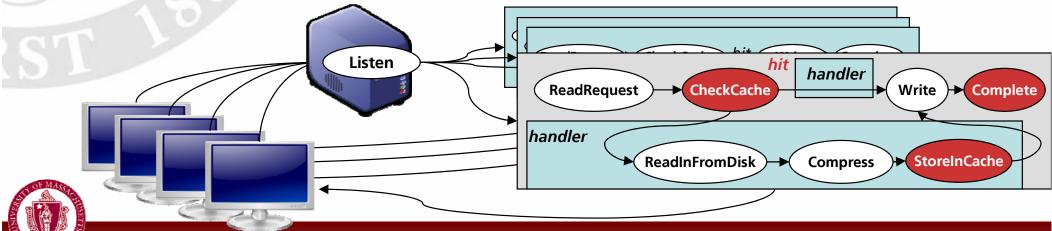


Supporting Concurrency

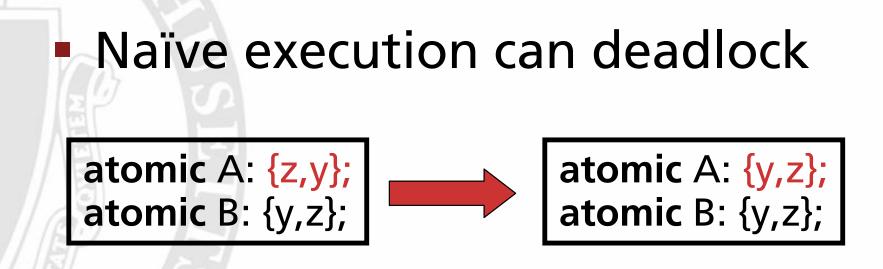
Many clients = concurrent flows

- Must keep cache consistent
- Atomicity constraints
 - Same name = mutual exclusion
 - Multiple names, reader/writer, per-client (see paper)

atomic CheckCache{cacheLock};atomic Complete{cacheLock};atomic StoreInCache{cacheLock};



Preventing Deadlock



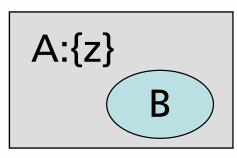
Establish canonical lock order Currently alphabetic by name



Preventing Deadlock, II

Harder with abstract nodes

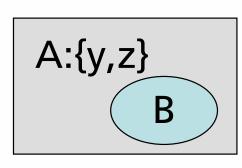
A = B; C = D; atomic A{z}; atomic B{y}; atomic C{y,z}

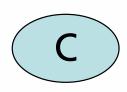




Solution: Elevate constraints

A = B; C = D; atomic A{y,z}; atomic B{y}; atomic C{y,z}





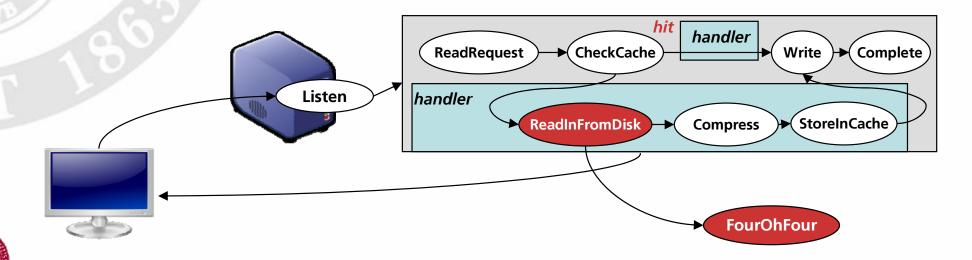
Finally: Handling Errors

What if image requested doesn't exist?

- Error = negative return value from component
- Can designate error handlers

Go to alternate paths on error

handle error ReadInFromDisk → FourOhFour;



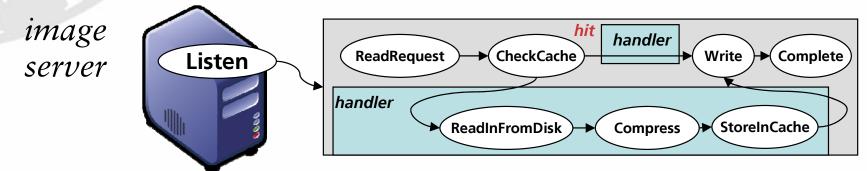
Complete Flux Image Server

```
source Listen → Image;
Image =
ReadRequest → CheckCache → Handler → Write → Complete;
Handler[_,_,hit] = ;
Handler[_,_,] = ReadFromDisk → Compress → StoreInCache;
atomic CheckCache: {cacheLock};
```

atomic CheckCache: {cacheLock}; atomic StoreInCache: {cacheLock}; atomic Complete: {cacheLock};

handle error ReadInFromDisk → FourOhFour;

- Concise, readable expression of server logic
 - No threads, etc.: simplifies programming, debugging



Outline

Intro to Flux: building a server Components, flow Atomicity, deadlock avoidance Performance results

- Server performance
- Performance prediction
- Future work



Results

Four servers:

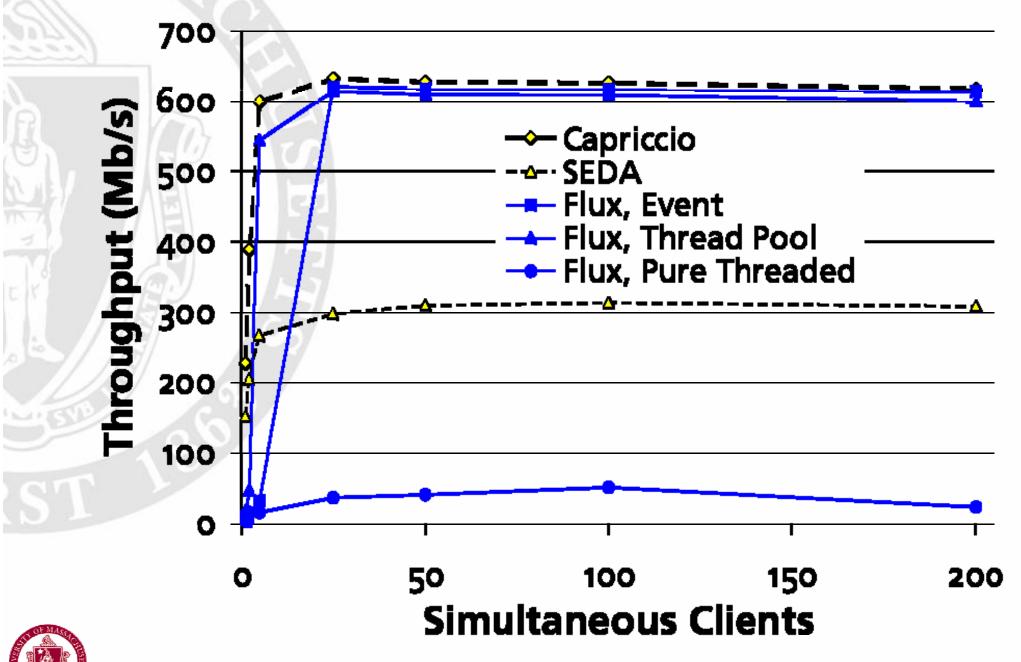
- Image server (+ libjpeg) [23 lines of Flux]
- Multi-player online game [54]
- BitTorrent (2 undergrads: 1 week!) [84]
- Web server (+ PHP) [36]

Evaluation

- Benchmark: variant of SPECweb99
- Three different runtimes
 - Thread, Thread pool, Event-Driven
- Compared to Capriccio [SOSP03], SEDA [SOSP01]

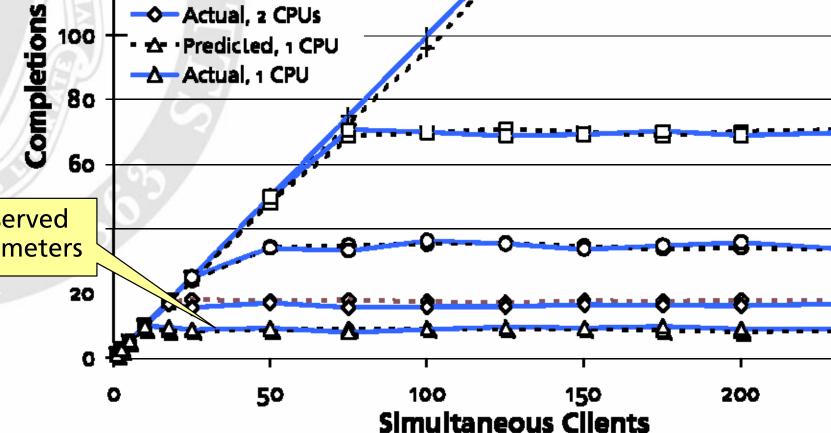


Web Server



Performance Prediction 180 +- · Predicted, 16 CPUs + Actual, 16 CPUs 160 - C - Predicted, 8 CPUs - C - Actual, 8 CPUs 140 - O · Predicted, 4 CPUs -O- Actual, 4 CPUs vn 120 Predicted, 2 CPUs - Actual, 2 CPUs 100 - 🛧 · Predicled, 1 CPU -Actual, 1 CPU 80

60 observed parameters 20



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Future Work

- Different runtimes
- Distributed architecturesClusters



- Embedded, power-aware systems
 - Turtles!
 - Embedded space similar to servers
 - eFlux: includes power constraints
 - Removes/adds flows dynamically in response to power



Conclusion

- Flux: language for server programming
 - Uses sequential code
 - Separates logic and runtime
 - Deadlock-free highperformance servers
 + simulators



- flux.cs.umass.edu
 - Hosted by Flux web server; download via Flux BitTorrent

flux: from Latin *fluxus,* p.p. of *fluere* = "to flow"



Backup Slides





Relaxed Atomicitv

Reader / writer constraints

Multiple readers or single writer

atomic ReadList: {listAccess?};
atomic AddToList: {listAccess!};

- Per-session constraints
 - One constraint per client / session

atomic AddHasChunk: {chunks(session)};

