# Probabilistic Context Free Grammars 

## Lecture \#16

Introduction to Natural Language Processing CMPSCI 585, Spring 2004<br><br>Andrew McCallum<br>(including slides from Jason Eisner)

## Ambiguity in Parsing

- Time flies like an arrow.
- Fruit flies like a banana.
- I saw the man with the telescope.


## How to solve this combinatorial explosion of ambiguity?

1. First try parsing without any weird rules, throwing them in only if needed.
2. Better: every rule has a weight. A tree's weight is total weight of all its rules.
Pick the overall "lightest" parse of sentence.
3. Can we pick the weights automatically? We'll get to this later ...

## CYK Parser

Input: A string of words, grammar in CNF
Output: yes/no
Data structure: $\mathrm{n} \times \mathrm{n}$ table
rows labeled 0 to $n-1$, columns 1 to $n$ cell (i,j) lists constituents spanning $\mathrm{i}, \mathrm{j}$

For each i from 1 to n
Add to ( $\mathrm{i}-1, \mathrm{i}$ ) all Nonterminals that could produce the word at (i-1,i)


## CYK Parser

## For width from 2 to n

For start from 0 to n-width
Define end to be start+width
For mid from start+1 to end-1
For every constituent in (start, mid)
For every constituent in (mid,end)
For all ways of combining them (if any)
Add the resulting constituent to (start,end).


$1 S \rightarrow N P$ VP
$6 \mathrm{~S} \rightarrow$ Vst NP
$2 S \rightarrow S$ PP
$1 \mathrm{VP} \rightarrow \mathrm{V}$ NP
$2 \mathrm{VP} \rightarrow \mathrm{VP}$ PP
$1 \mathrm{NP} \rightarrow$ Det N
2 NP $\rightarrow$ NP PP
3 NP $\rightarrow$ NP NP
$0 \mathrm{PP} \rightarrow \mathrm{PNP}$


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2 NP $\rightarrow$ NP PP
$3 N P \rightarrow N P N P$
$0 \mathrm{PP} \rightarrow \mathrm{PNP}$









time 1 flies 2 like 3 an 4 arrow 5

$1 S \rightarrow N P$ VP
$6 S \rightarrow$ Vst NP
$2 S \rightarrow S P P$
$1 \mathrm{VP} \rightarrow \mathrm{V}$ NP
$2 \mathrm{VP} \rightarrow \mathrm{VP}$ PP
$1 \mathrm{NP} \rightarrow$ Det $N$
2 NP $\rightarrow$ NP PP
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time 1 flies 2 like 3 an 4 arrow 5

|  | $\begin{array}{ll} \hline \text { NP } & 3 \\ \text { Vst } & 3 \end{array}$ | $\begin{array}{ll}\text { NP } & 10 \\ S & 8 \\ S & 13\end{array}$ |  |  | $\begin{array}{ll}\text { NP } & 24 \\ \text { S } & 22\end{array}$ |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  |  |  |  | $\begin{aligned} & 6 S \rightarrow \text { Vst NP } \\ & 2 S \rightarrow S \text { PP } \end{aligned}$ |
| 1 |  | $\begin{array}{ll}\text { NP } & 4 \\ \text { VP } & 4\end{array}$ |  |  | $\begin{array}{ll} \mathrm{NP} & 18 \\ \mathrm{~S} & 21 \\ \mathrm{VP} & 18 \end{array}$ | $\begin{aligned} & 1 \mathrm{VP} \rightarrow \mathrm{~V} \mathrm{NP} \\ & 2 \mathrm{VP} \rightarrow \mathrm{VP} \mathrm{PP} \end{aligned}$ |
| 2 |  |  | P 2 V 5 |  | $\begin{array}{\|ll} \hline \text { PP } & 12 \\ \text { VP } & 16 \end{array}$ | $\begin{aligned} & 1 \mathrm{NP} \rightarrow \operatorname{Det} \mathrm{~N} \\ & 2 \mathrm{NP} \rightarrow \mathrm{NP} \mathrm{PP} \end{aligned}$ |
| 3 |  |  |  | Det 1 | NP 10 | 3 NP $\rightarrow$ NP NP |
| 4 |  |  |  |  | N 8 | $0 \mathrm{PP} \rightarrow \mathrm{PNP}$ |



|  | e 1 | 2 lik | 3 | 4 ar | 5 |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  | $\begin{array}{\|ll\|} \hline \text { NP } & 3 \\ \text { Vst } & 3 \end{array}$ | $\begin{array}{ll} \mathrm{NP} & 10 \\ \mathrm{~S} & 8 \\ \mathrm{~S} & 13 \end{array}$ |  |  | $N P$ 24 <br> $S$ 22 <br> $S$ 27 |
|  |  | $\begin{array}{ll}\text { NP } & 4 \\ \text { VP } & 4\end{array}$ |  |  | NP 18 <br> S 21 <br> VP 18 |
|  |  |  | P 2 V 5 |  | $\begin{array}{ll} \hline \text { PP } & 12 \\ \text { VP } & 16 \end{array}$ |
|  |  |  |  | Det 1 | NP 10 |
|  |  |  |  |  | N 8 |

$1 S \rightarrow N P$ VP
$6 S \rightarrow$ Vst NP
$2 S \rightarrow S P P$
$1 \mathrm{VP} \rightarrow \mathrm{V}$ NP
$2 \mathrm{VP} \rightarrow \mathrm{VP}$ PP
$1 \mathrm{NP} \rightarrow$ Det N
$2 N P \rightarrow N P$ PP
$3 N P \rightarrow N P N P$
$0 \mathrm{PP} \rightarrow \mathrm{PNP}$


$1 \mathrm{~S} \rightarrow \mathrm{NP}$ VP
$6 \mathrm{~S} \rightarrow \mathrm{Vst} \mathrm{NP}$
$2 \mathrm{~S} \rightarrow \mathrm{~S} P \mathrm{PP}$
$1 \mathrm{VP} \rightarrow \mathrm{V} \mathrm{NP}$
$2 \mathrm{VP} \rightarrow \mathrm{VP} \mathrm{PP}$
$1 \mathrm{NP} \rightarrow$ Det N
$2 \mathrm{NP} \rightarrow \mathrm{NP} \mathrm{PP}$
$3 \mathrm{NP} \rightarrow \mathrm{NP} N P$
$0 \mathrm{PP} \rightarrow \mathrm{PNP}$
time 1 flies 2 like 3 an 4 arrow 5

$1 S \rightarrow N P$ VP
$6 \mathrm{~S} \rightarrow$ Vst NP
$2 S \rightarrow$ SPP
$1 \mathrm{VP} \rightarrow \mathrm{VNP}$
$2 \mathrm{VP} \rightarrow \mathrm{VP} P \mathrm{P}$
$1 \mathrm{NP} \rightarrow$ Det N
$2 \mathrm{NP} \rightarrow \mathrm{NP}$ PP
$3 \mathrm{NP} \rightarrow \mathrm{NP}$ NP
$0 \mathrm{PP} \rightarrow \mathrm{PNP}$

$1 S \rightarrow N P V P$
$6 S \rightarrow$ Vst NP
$2 S \rightarrow$ S PP
$1 \mathrm{VP} \rightarrow \mathrm{VNP}$
$2 \mathrm{VP} \rightarrow \mathrm{VP} P \mathrm{P}$
$1 \mathrm{NP} \rightarrow \operatorname{Det} \mathrm{N}$
$2 N P \rightarrow N P$ PP
$3 \mathrm{NP} \rightarrow \mathrm{NPNP}$
$0 \mathrm{PP} \rightarrow \mathrm{PNP}$



|  | me 1 fl | S | lik | 3 | 4 | 5 |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | NP 3 <br> Vst 3 | NP | $\begin{aligned} & \hline 10 \\ & 8 \\ & 13 \end{aligned}$ |  |  | $N P$ 24 <br> $S$ 22 <br> $S$ 27 <br> $N P$ 24 <br> $S$ 27 <br> $S$ 22 <br> $S$ 27 |  |
| 1 |  |  |  |  |  | NP 18 <br> S 21 <br> VP 18 | $\begin{aligned} & 6 \mathrm{~S} \rightarrow \mathrm{Vst} \mathrm{NP} \\ & 2 \mathrm{~S} \rightarrow \mathrm{SPP} \\ & 1 \mathrm{VP} \rightarrow \mathrm{VNP} \end{aligned}$ |
| 2 |  |  |  | $\begin{array}{ll}\text { P } 2 \\ \text { V } & 5\end{array}$ |  | PP 12 <br> VP 16 | $\begin{aligned} & 2 \mathrm{VP} \rightarrow \mathrm{VP} \mathrm{PP} \\ & 1 \mathrm{NP} \rightarrow \operatorname{Det} \mathrm{~N} \end{aligned}$ |
| 3 |  |  |  |  | Det 1 | NP 10 | $3 \mathrm{NP} \rightarrow \mathrm{NP} \mathrm{NP}$ |
| 4 |  |  |  |  |  | N 8 | $0 \mathrm{PP} \rightarrow \mathrm{PNP}$ |



| time 1 flies |  |  | lik | 3 | 4 | arr | 5 |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | NP 3 <br> Vst 3 | NP S S | $\begin{aligned} & 10 \\ & 8 \\ & 13 \end{aligned}$ |  |  |  | NP 24 <br> $S$ 22 <br> $S$ 27 <br> NP 24 <br> $S$ 27 <br> $S$ 22 <br> $S$ 27 |  |
| 1 |  |  | 4 |  |  |  | $\begin{array}{lll}\text { NP } & 18 \\ \mathrm{~S} & 21 \\ \mathrm{VP} & 18\end{array}$ | $\begin{aligned} & 1 \mathrm{~S} \rightarrow \mathrm{NP} \text { VP } \\ & 6 \mathrm{~S} \rightarrow \mathrm{Vst} \mathrm{NP} \\ & 2 \mathrm{~S} \rightarrow \mathrm{SPP} \\ & 1 \mathrm{VP} \rightarrow \mathrm{VNP} \end{aligned}$ |
| 2 |  |  |  | $\begin{array}{ll}\text { P } & 2 \\ \mathrm{~V} & 5\end{array}$ |  |  | $P P$ 12 <br> VP 16 | $\begin{aligned} & 2 \mathrm{VP} \rightarrow \mathrm{VP} \mathrm{PP} \\ & 1 \mathrm{NP} \rightarrow \operatorname{Det} \mathrm{~N} \end{aligned}$ |
| 3 |  |  |  |  | Det | 1 | NP 10 | $\begin{aligned} & 2 \mathrm{NP} \rightarrow \mathrm{NP} \mathrm{PP} \\ & 3 \mathrm{NP} \rightarrow \mathrm{NP} \mathrm{NP} \end{aligned}$ |
| 4 |  |  |  |  |  |  | N 8 | $0 \mathrm{PP} \rightarrow \mathrm{PNP}$ |

## Which entries do we need?

time 1 flies 2 like 3 an 4 arrow 5


$$
\begin{aligned}
& 1 \mathrm{~S} \rightarrow \mathrm{NP} \mathrm{VP} \\
& 6 \mathrm{~S} \rightarrow \mathrm{Vst} N P \\
& 2 \mathrm{~S} \rightarrow \mathrm{SPP} \\
& 1 \mathrm{VP} \rightarrow \mathrm{~V} \text { NP } \\
& 2 \mathrm{VP} \rightarrow \mathrm{VP} P \mathrm{PP} \\
& 1 \mathrm{NP} \rightarrow \operatorname{Det} \mathrm{~N} \\
& 2 \mathrm{NP} \rightarrow \mathrm{NP} \mathrm{PP} \\
& 3 \mathrm{NP} \rightarrow \mathrm{NP} \mathrm{NP} \\
& 0 \mathrm{PP} \rightarrow \mathrm{PNP}
\end{aligned}
$$

## Which entries do we need?

time 1 flies 2 like 3 an 4 arrow 5

| 0 | $\begin{array}{ll} \hline \text { NP } & 3 \\ \text { Vst } & 3 \end{array}$ |  |  |  | NP 24 <br> S 22 <br> S 27 <br> NP 24 <br> S 27 <br> S 22 <br> S 27 |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 1 |  | $\begin{array}{\|ll\|} \hline \text { NP } & 4 \\ \text { VP } & 4 \end{array}$ |  |  | $\begin{array}{ll} \hline N P & 18 \\ S & 21 \\ \text { VP } & 18 \end{array}$ |
| 2 |  |  | P 2 V 5 |  | $\begin{array}{\|ll} \hline \text { PP } & 12 \\ \text { VP } & 16 \end{array}$ |
| 3 |  |  |  | Det 1 | NP 10 |
| 4 |  |  |  |  | N 8 |

$1 \mathrm{~S} \rightarrow \mathrm{NP}$ VP
$6 \mathrm{~S} \rightarrow$ Vst NP
$2 S \rightarrow S P P$
$1 \mathrm{VP} \rightarrow \mathrm{VNP}$
$2 \mathrm{VP} \rightarrow \mathrm{VP}$ PP
$1 \mathrm{NP} \rightarrow \operatorname{Det} \mathrm{N}$
$2 N P \rightarrow N P$ PP
$3 \mathrm{NP} \rightarrow \mathrm{NP}$ NP
$0 \mathrm{PP} \rightarrow \mathrm{PNP}$


## since it just breeds worse options

| ne 1 flies 2 lik |  |  | 3 | 5 |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | $\begin{array}{\|ll\|} \hline \text { NP } & 3 \\ \text { Vst } & 3 \end{array}$ | $\begin{array}{ll} \hline N P & 10 \\ S & 8 \\ \hline S & 13 \\ \hline \end{array}$ |  |  | $N P$ 24 <br> $S$ 22 <br> $S$ 27 <br> $N P$ 24 <br> $S$ 27 <br> $S$ 22 <br> $S$ 27 |
| 1 |  | $\begin{array}{\|ll\|} \hline \text { NP } & 4 \\ \text { VP } & 4 \end{array}$ |  |  | $\begin{array}{ll} \hline N P & 18 \\ S & 21 \\ \text { VP } & 18 \end{array}$ |
| 2 |  |  | P 2 V 5 |  | $\begin{array}{ll} \hline \mathrm{PP} & 12 \\ \mathrm{VP} & 16 \end{array}$ |
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$1 \mathrm{~S} \rightarrow \mathrm{NP}$ VP
$6 \mathrm{~S} \rightarrow \mathrm{Vst} \mathrm{NP}$
$2 S \rightarrow S P P$
$1 \mathrm{VP} \rightarrow \mathrm{VNP}$
$2 \mathrm{VP} \rightarrow \mathrm{VP}$ PP
$1 \mathrm{NP} \rightarrow \operatorname{Det} \mathrm{N}$
$2 N P \rightarrow N P$ PP
$3 \mathrm{NP} \rightarrow \mathrm{NP}$ NP
$0 \mathrm{PP} \rightarrow \mathrm{PNP}$

Keep only best-in-class!


Keep only best-in-class!
(and backpointers so you can recover parse)

| time 1 flies 2 like 3 an 4 arrow 5 |
| :--- |
| NP 3 NP 10    <br> Vst 3 S 8    |


|  | $S \rightarrow$ NP VP |
| :---: | :---: |
| 6 | $\mathrm{S} \rightarrow$ Vst NP |
| 2 | $S \rightarrow$ SPP |
| 1 | $\mathrm{VP} \rightarrow \mathrm{VNP}$ |
| 2 | $\mathrm{VP} \rightarrow \mathrm{VP}$ PP |
| 1 | $N P \rightarrow \operatorname{Det} N$ |
| 2 | $N P \rightarrow$ PP PP |
| 3 | $N P \rightarrow N P N P$ |
|  | $\mathrm{PP} \rightarrow \mathrm{PNP}$ |

## Probabilistic Trees

- Instead of lightest weight tree, take highest probability tree
- Given any tree, your assignment generator would have some probability of producing it!
- Just like using n-grams to choose among strings ...
- What is the probability of this tree?



## Probabilistic Trees

- Instead of lightest weight tree, take highest probability tree
- Given any tree, your assignment generator would have some probability of producing it!
- Just like using n-grams to choose among strings ...
- What is the probability of this tree?
- 
- You rolled a lot of independent dice...



## Chain rule: One word at a time

p (time flies like an arrow)
$=p$ (time)

* $p$ (flies | time)
* p(like | time flies)
* p(an | time flies like)
* p(arrow | time flies like an)


## Chain rule + backoff (to get trigram model)

p (time flies like an arrow)
$=p($ time $)$

* p(flies | time)
* p(like | time flies)
* p(an | time flies like)
* $p$ (arrow | time flies like an)


## Chain rule - written differently

p (time flies like an arrow)
$=\begin{aligned} & \mathrm{p} \text { (time) } \\ & * \mathrm{p} \text { (time flies | time) }\end{aligned}$

* p(time flies like | time flies)
* p(time flies like an | time flies like)
* p(time flies like an arrow | time flies like an)

Proof: $p(x, y \mid x)=p(x \mid x) * p(y \mid x, x)=1 * p(y \mid x)$

## Chain rule + backoff

$p$ (time flies like an arrow)

$$
=\underset{*}{p(\text { time }} \text { p(time flies I time) }
$$

* p(time flies like I time flies)
* p(time flies like an I time flies like)
* p (time flies like an arrow I time flies like an)

Proof: $p(x, y \mid x)=p(x \mid x)^{*} p(y \mid x, x)=1^{*} p(y \mid x)$

## Chain rule: One node at a time



Chain rule + backoff


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Simplified notation


Already have a CKY alg for weights


Just let $\mathbf{w}(x \rightarrow y z)=-\log \mathbf{p}(x \rightarrow y z \mid x)$ Then lightest tree has highest prob ${ }^{49}$
time 1 flies 2 like 3 an 4 arrow 5


$$
\begin{array}{lll} 
& & \\
& & 2^{-2} \\
1 & \mathrm{~S} \rightarrow \mathrm{NP} \mathrm{VP} \\
6 \mathrm{~S} \rightarrow \mathrm{Vst} \mathrm{NP} \\
2 & \mathrm{~S} \rightarrow \mathrm{~S} \mathrm{PP} \\
1 & \mathrm{VP} \rightarrow \mathrm{~V} \mathrm{NP} \\
2 \mathrm{VP} \rightarrow \mathrm{VP} \mathrm{PP} \\
1 \mathrm{NP} \rightarrow \mathrm{Det} \mathrm{~N} \\
2 \mathrm{NP} \rightarrow \mathrm{NP} \mathrm{PP} \\
3 \mathrm{NP} \rightarrow \mathrm{NP} \mathrm{NP} \\
0 \mathrm{PP} \rightarrow \mathrm{PNP}
\end{array}
$$

Need only best-in-class to get best parse



## Why probabilities not weights?

- We just saw probabilities are really just a special case of weights ...
- ... but we can estimate them from training data by counting and smoothing! Use all of our lovely probability theory machinery!


## Probabilistic Context Free Grammars <br> (PCFGs)

A PCFG $G$ consists of the usual parts of a CFG

- A set of terminals, $\left\{w^{k}\right\}, k=1, \ldots, V$

■ A set of nonterminals, $\left\{N^{i}\right\}, i=1, \ldots, n$

- A designated start symbol, $N^{1}$

■ A set of rules, $\left\{N^{i} \rightarrow \zeta^{j}\right\}$, (where $\zeta^{j}$ is a sequence of terminals and nonterminals)
and

- A corresponding set of probabilities on rules such that:

$$
\forall i \quad \sum_{j} P\left(N^{i} \rightarrow \zeta^{j}\right)=1
$$

## PCFG notation

Sentence: sequence of words $w_{1} \cdots w_{m}$
$w_{a b}$ : the subsequence $w_{a} \cdots w_{b}$
$N_{a b}^{i}$ : nonterminal $N^{i}$ dominates $w_{a} \cdots w_{b}$

$N^{i} \stackrel{*}{\Rightarrow} \zeta:$ Repeated derivation from $N^{i}$ gives $\zeta$.

## PCFG probability of a string

$$
\begin{aligned}
P\left(w_{1 n}\right) & =\sum_{t} P\left(w_{1 n}, t\right) \quad t \text { a parse of } w_{1 n} \\
& =\sum_{\left\{t: y i e l d(t)=w_{1 n}\right\}} P(t)
\end{aligned}
$$

## A simple PCFG (in CNF)

| $\mathrm{S} \rightarrow \mathrm{NP} \mathrm{VP}$ | 1.0 | $\mathrm{NP} \rightarrow$ NP PP | 0.4 |
| :--- | :--- | :--- | :--- |
| $\mathrm{PP} \rightarrow \mathrm{P} \mathrm{NP}$ | 1.0 | $\mathrm{NP} \rightarrow$ astronomers | 0.1 |
| $\mathrm{VP} \rightarrow \mathrm{V} \mathrm{NP}$ | 0.7 | $\mathrm{NP} \rightarrow$ ears | 0.18 |
| $\mathrm{VP} \rightarrow \mathrm{VP} \mathrm{PP}$ | 0.3 | $\mathrm{NP} \rightarrow$ saw | 0.04 |
| $\mathrm{P} \rightarrow$ with | 1.0 | $\mathrm{NP} \rightarrow$ stars | 0.18 |
| $\mathrm{~V} \rightarrow$ saw | 1.0 | $\mathrm{NP} \rightarrow$ telescopes | 0.1 |



## The two parse trees' probabilities and the sentence probability

$$
\begin{aligned}
P\left(t_{1}\right)= & 1.0 \times 0.1 \times 0.7 \times 1.0 \times 0.4 \\
& \times 0.18 \times 1.0 \times 1.0 \times 0.18 \\
= & 0.0009072 \\
P\left(t_{2}\right)= & 1.0 \times 0.1 \times 0.3 \times 0.7 \times 1.0 \\
& \times 0.18 \times 1.0 \times 1.0 \times 0.18 \\
= & 0.0006804 \\
P\left(w_{15}\right)= & P\left(t_{1}\right)+P\left(t_{2}\right)=0.0015876
\end{aligned}
$$

## Assumptions of PCFGs

1. Place invariance (like time invariance in HMM):

$$
\forall k \quad P\left(N_{k(k+c)}^{j} \rightarrow \zeta\right) \text { is the same }
$$

2. Context-free:

$$
P\left(N_{k l}^{j} \rightarrow \zeta \mid \text { words outside } w_{k} \ldots w_{l}\right)=P\left(N_{k l}^{j} \rightarrow \zeta\right)
$$

3. Ancestor-free:

$$
P\left(N_{k l}^{j} \rightarrow \zeta \mid \text { ancestor nodes of } N_{k l}^{j}\right)=P\left(N_{k l}^{j} \rightarrow \zeta\right)
$$

The sufficient statistics of a PCFG are thus simply counts of how often different local tree configurations occurred (= counts of which grammar rules were applied).

## Some features of PCFGs

Reasons to use a PCFG, and some idea of their limitations:

- Partial solution for grammar ambiguity: a PCFG gives some idea of the plausibility of a sentence.
- But, in the simple case, not a very good idea, as independence assumptions are two strong (e.g., not lexicalized).
- Gives a probabilistic language model for English.
- In the simple case, a PCFG is a worse language model for English than a trigram model.

■ Better for grammar induction (Gold 1967 vs. Horning 1969)

■ Robustness. (Admit everything with low probability.)

## Some features of PCFGs

- A PCFG encodes certain biases, e.g., that smaller trees are normally more probable.
- One can hope to combine the strengths of a PCFG and a trigram model.

We'll look at simple PCFGs first. They have certain inadequacies, but we'll see that most of the state-of-the-art probabilistic parsers are fundamentally PCFG models, just with various enrichments to the grammar

## A slightly different task

- Been asking: What is probability of generating a given tree?
- To pick tree with highest prob: useful in parsing.
- But could also ask: What is probability of generating a given string with the generator?
- To pick string with highest prob: useful in speech recognition, as substitute for an n-gram model.
- ("Put the file in the folder" vs. "Put the file and the folder")
- To get prob of generating string, must add up probabilities of all trees for the string ...


## Could just add up the parse probabilities



## Any more efficient way?



## Add as we go ... (the "inside algorithm")



Add as we go ... (the "inside algorithm")


## Inside and Outside Probabilities

Probability of all possible rule re-writes for generating words inside position p to q , given that non-terminal $j$ exactly spans $p$ to $q$.

$$
\text { Inside }=\beta_{j}(p, q)=P\left(w_{p q} \mid N_{p q}^{j}, G\right)
$$

Probability of all possible rule re-writes for generating words outside position p to q , and that non-terminal $j$ exactly spans $p$ to $q$.

$$
\text { Outside }=\alpha_{j}(p, q)=P\left(w_{1(p-1)}, N_{p q}^{j}, w_{(q+1) m} \mid G\right)
$$

## Inside and Outside Probabilities



## Inside \& Outside Probabilities



$$
=p(\mathrm{VP}(1,5) \mid \text { time flies like an arrow today, } \mathrm{S})
$$

So $\alpha \vee p(1,5) * \beta \vee p(1,5) / \beta_{\mathbf{s}}(0,6)$
is probability that there is a VP here, given all of the observed data (words)

## Probability of a string

## Inside probability

$$
\begin{aligned}
P\left(w_{1 m} \mid G\right) & =P\left(N^{1} \Rightarrow w_{1 m} \mid G\right) \\
& =P\left(w_{1 m}, N_{1 m}^{1}, G\right)=\beta_{1}(1, m)
\end{aligned}
$$

Base case: We want to find $\beta_{j}(k, k)$ (the probability of a rule $\left.N^{j} \rightarrow w_{k}\right):$

$$
\begin{aligned}
\beta_{j}(k, k) & =P\left(w_{k} \mid N_{k k}^{j}, G\right) \\
& =P\left(N^{j} \rightarrow w_{k} \mid G\right)
\end{aligned}
$$

## Probability of a string

Induction: We want to find $\beta_{j}(p, q)$, for $p<q$. As this is the inductive step using a Chomsky Normal Form grammar, the first rule must be of the form $N^{j} \rightarrow N^{r} \quad N^{S}$, so we can proceed by induction, dividing the string in two in various places and summing the result:


These inside probabilities can be calculated bottom up.

For all $j$,

$$
\begin{aligned}
\beta_{j}(p, q)= & P\left(w_{p q} \mid N_{p q}^{j}, G\right) \\
= & \sum_{r, s}^{q-1} \sum_{d=p}^{q-1} P\left(w_{p d}, N_{p d}^{r}, w_{(d+1) q}, N_{(d+1) q}^{s} \mid N_{p q}^{j}, G\right) \\
= & \sum_{r, s} \sum_{d=p}^{q-1} P\left(N_{p d}^{r}, N_{(d+1) q}^{s} \mid N_{p q}^{j}, G\right) \\
& P\left(w_{p d} \mid N_{p q}^{j}, N_{p d}^{r}, N_{(d+1) q}^{s}, G\right) \\
& P\left(w_{(d+1) q} \mid N_{p q}^{j}, N_{p d}^{r}, N_{(d+1) q}^{s}, w_{p d}, G\right) \\
= & \sum_{r, s} \sum_{d=p}^{q-1} P\left(N_{p d}^{r}, N_{(d+1) q}^{s} \mid N_{p q}^{j}, G\right) \\
& P\left(w_{p d} \mid N_{p d}^{r}, G\right) P\left(w_{(d+1) q} \mid N_{(d+1) q}^{s}, G\right) \\
= & \sum_{r, s} \sum_{d=p}^{q-1} P\left(N^{j} \rightarrow N^{r} N^{s}\right) \beta_{r}(p, d) \beta_{S}(d+1, q)
\end{aligned}
$$

## Inside probabilities as CYK

|  | 1 | 2 | 3 | 4 | 5 |  |
| :--- | :--- | :--- | :--- | :--- | :--- | :---: |
| 1 | $\beta_{\mathrm{NP}}=0.1$ |  | $\beta_{\mathrm{S}}=0.0126$ |  | $\beta_{\mathrm{S}}=0.0015876$ |  |
| 2 |  | $\beta_{\mathrm{NP}}=0.04$ <br> $\beta_{\mathrm{V}}=1.0$ | $\beta_{\mathrm{VP}}=0.126$ |  | $\beta_{\mathrm{VP}}=0.015876$ |  |
| 3 |  |  | $\beta_{\mathrm{NP}}=0.18$ |  | $\beta_{\mathrm{NP}}=0.01296$ |  |
| 4 |  |  |  | $\beta_{\mathrm{P}}=1.0$ | $\beta_{\mathrm{PP}}=0.18$ |  |
| 5 |  |  |  | $\beta_{\mathrm{NP}}=0.18$ |  |  |
|  | astronomers | saw | stars | with | ears |  |

## Outside probabilities

Probability of a string: For any $k, 1 \leq k \leq m$,

$$
\begin{aligned}
P\left(w_{1 m} \mid G\right)= & \sum_{j} P\left(w_{1(k-1)}, w_{k}, w_{(k+1) m}, N_{k k}^{j} \mid G\right) \\
= & \sum_{j} P\left(w_{1(k-1)}, N_{k k}^{j}, w_{(k+1) m} \mid G\right) \\
& \times P\left(w_{k} \mid w_{1(k-1)}, N_{k k}^{j}, w_{(k+1) n}, G\right) \\
= & \sum_{j} \alpha_{j}(k, k) P\left(N^{j} \rightarrow w_{k}\right)
\end{aligned}
$$

Inductive (DP) calculation: One calculates the outside probabilities top down (after determining the inside probabilities).

## Outside probabilities

## Base Case:

$$
\begin{aligned}
& \alpha_{1}(1, m)=1 \\
& \alpha_{j}(1, m)=0, \text { for } j \neq 1
\end{aligned}
$$

Inductive Case: it's either a left or right branch - we will some over both possibilities and calculate using outside and inside probabilities


## Outside probabilities, Inductive case

$$
\alpha_{j}(p, q)=\left[\sum_{f, g} \sum_{e=q+1}^{m} P\left(w_{1(p-1)}, w_{(q+1) m}, N_{p e}^{f}, N_{p q}^{j}, N_{(q+1) e}^{g}\right)\right]
$$

$$
+\left[\sum_{f, g}^{p-1} \sum_{e=1}^{p} P\left(w_{1(p-1)}, w_{(q+1) m}, N_{e q}^{f}, N_{e(p-1)}^{g}, N_{p q}^{j}\right)\right]
$$

$$
=\left[\sum_{f, g} \sum_{e=q+1}^{m} P\left(w_{1(p-1)}, w_{(e+1) m}, N_{p e}^{f}\right) P\left(N_{p q}^{j}, N_{(q+1) e}^{g} \mid N_{p e}^{f}\right)\right.
$$

$$
\left.\times P\left(w_{(q+1) e} \mid N_{(q+1) e}^{g}\right)\right]+\left[\sum_{f, g}^{p-1} \sum_{e=1}^{p} P\left(w_{1(e-1)}, w_{(q+1) m}, N_{e q}^{f}\right)\right.
$$

$$
\left.\times P\left(N_{e(p-1)}^{g}, N_{p q}^{j} \mid N_{e q}^{f}\right) P\left(w_{e(p-1)} \mid N_{e(p-1}^{g}\right)\right]
$$

$$
=\left[\sum_{f, g} \sum_{e=q+1}^{m} \alpha_{f}(p, e) P\left(N^{f} \rightarrow N^{j} N^{g}\right) \beta_{g}(q+1, e)\right]
$$

$$
+\left[\sum_{f, g}^{p-1} \sum_{e=1}^{p-1} \alpha_{f}(e, q) P\left(N^{f} \rightarrow N^{g} N^{j}\right) \beta_{g}(e, p-1)\right]
$$

## Probability that a rule is used

As with a HMM, we can form a product of the inside and outside probabilities. This time:

$$
\begin{aligned}
& \quad \alpha_{j}(p, q) \beta_{j}(p, q) \\
& =P\left(w_{1(p-1)}, N_{p q}^{j}, w_{(q+1) m} \mid G\right) P\left(w_{p q} \mid N_{p q}^{j}, G\right) \\
& =P\left(w_{1 m}, N_{p q}^{j} \mid G\right) \\
& P\left(N_{p q}^{j} \mid w_{1 m}, G\right)=\frac{P\left(N_{p q}^{j} \mid w_{1 m}, G\right)}{P\left(w_{1 m} \mid G\right)}=\frac{\alpha_{j}(p, q) \beta_{j}(p, q)}{\beta_{1}(1, m)}
\end{aligned}
$$

This is an "expected count" for the number of times this rule occurred.

## Overall probability of a node existing

As with a HMM, we can form a product of the inside and outside probabilities. This time:

$$
\begin{aligned}
& \alpha_{j}(p, q) \beta_{j}(p, q) \\
& =P\left(w_{1(p-1)}, N_{p q}^{j}, w_{(q+1) m} \mid G\right) P\left(w_{p q} \mid N_{p q}^{j}, G\right) \\
& =P\left(w_{1 m}, N_{p q}^{j} \mid G\right)
\end{aligned}
$$

Therefore,

$$
p\left(w_{1 m}, N_{p q} \mid G\right)=\sum_{j} \alpha_{j}(p, q) \beta_{j}(p, q)
$$

Just in the cases of the root node and the preterminals, we know there will always be some such constituent.

## Learning PCFGs (1)

- We would like to calculate how often each rule is used:

$$
\hat{P}\left(N^{j} \rightarrow \zeta\right)=\frac{C\left(N^{j} \rightarrow \zeta\right)}{\sum_{\gamma} C\left(N^{j} \rightarrow \gamma\right)}
$$

- If we have labeled data, we count and find out
- Relative frequency again gives maximum likelihood probability estimates
- This is the motivation for building Treebanks of handparsed sentences


## Learning PCFGs (2) Inside-Outside

- Otherwise we work iteratively from expectations of current model.
- We construct an EM training algorithm, as for HMMs
- For each sentence, at each iteration, we work out expectation of how often each rule is used using inside and outside probabilities
- We assume sentences are independent and sum expectations over parses of each
- We re-estimate rules based on these 'counts'

