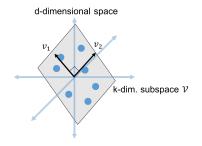
# COMPSCI 514: ALGORITHMS FOR DATA SCIENCE

Andrew McGregor

Lecture 14

#### LAST CLASS: EMBEDDING WITH ASSUMPTIONS

**Set Up:** Assume that data points  $\vec{x_1}, \ldots, \vec{x_n} \in \mathbb{R}^d$  lie in some *k*-dimensional subspace  $\mathcal{V}$  of  $\mathbb{R}^d$ .



Let  $\vec{v}_1, \ldots, \vec{v}_k$  be an orthonormal basis for  $\mathcal{V}$  and  $\mathbf{V} \in \mathbb{R}^{d \times k}$  be the matrix with these vectors as its columns.

$$\|\mathbf{V}^T \vec{x}_i - \mathbf{V}^T \vec{x}_j\|_2^2 = \|\vec{x}_i - \vec{x}_j\|_2^2.$$

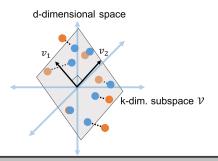
Letting  $\tilde{x}_i = \mathbf{V}^T \vec{x}_i$ , we have a perfect embedding from  $\mathcal{V}$  into  $\mathbb{R}^k$ .

# LAST CLASS: PROJECTION VIEW

**Warm-Up:** If  $\vec{x_1}, \ldots, \vec{x_n}$  lie in a *k*-dimensional subspace  $\mathcal{V}$  with orthonormal basis  $\mathbf{V} \in \mathbb{R}^{d \times k}$ , the data matrix can be written as

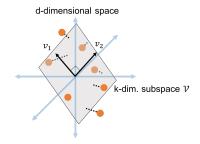
 $\mathbf{X} = \mathbf{X}\mathbf{V}\mathbf{V}^{\mathsf{T}} = \mathbf{C}\mathbf{V}^{\mathsf{T}}$ 

• **VV**<sup>T</sup> is a projection matrix, which projects the rows of **X** (the data points  $\vec{x}_1, \ldots, \vec{x}_n$ ) onto the subspace  $\mathcal{V}$ .



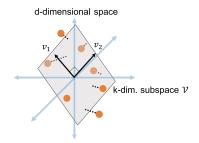
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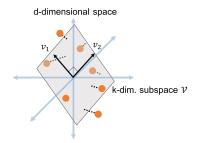


Letting  $\vec{v}_1, \ldots, \vec{v}_k$  be an orthonormal basis for  $\mathcal{V}$  and  $\mathbf{V} \in \mathbb{R}^{d \times k}$  be the matrix with these vectors as its columns,  $\mathbf{V}^T \vec{x}_i \in \mathbb{R}^k$  is still a good embedding for  $\mathbf{x}_i \in \mathbb{R}^d$  and  $\mathbf{X}\mathbf{V}\mathbf{V}^T$  is still a good approximation for  $\mathbf{X}$  in the sense:

$$\mathbf{X}\mathbf{V}\mathbf{V}^{\mathsf{T}} = \underset{\mathbf{B} \text{ with rows in } \mathcal{V}}{\arg\min} \|\mathbf{X} - \mathbf{B}\|_{F}^{2}$$

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Will show first show this and then investigate how do we find  $\mathcal{V}$  and  $\mathbf{V}$ ?

**Exercise 1 :** Show that  $\mathbf{VV}^{T}$  is idempotent. I.e., for any  $\vec{y} \in \mathbb{R}^{d}$ :

$$(\mathbf{V}\mathbf{V}^{\mathsf{T}})(\mathbf{V}\mathbf{V}^{\mathsf{T}})\vec{y} = (\mathbf{V}\mathbf{V}^{\mathsf{T}})\vec{y}$$

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**Exercise 2:** The projection is orthogonal to its complement: For any  $\vec{y} \in \mathbb{R}^d$ ,  $\langle \mathbf{V} \mathbf{V}^T \vec{y}, \vec{y} - \mathbf{V} \mathbf{V}^T \vec{y} \rangle = 0$ 

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**Implies the Pythagorean Theorem:** Show that for any  $\vec{y} \in \mathbb{R}^d$ ,

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Follows since  $\vec{y} = (\vec{y} - (\mathbf{V}\mathbf{V}^T)\vec{y}) + (\mathbf{V}\mathbf{V}^T)\vec{y}$  and

$$\|\vec{a} + \vec{b}\|_2^2 = \|\vec{a}\|_2^2 + \|\vec{b}\|_2^2 + 2\langle \vec{a}, \vec{b} \rangle$$

Let  $\mathbf{V} \in \mathbb{R}^{n \times k}$  have orthonormal columns and let  $\vec{y} \in \mathbb{R}^n$ . Then the Pythagorean Theorem proves that  $\mathbf{V}\mathbf{V}^T\vec{y}$  is the closest vector to  $\vec{y}$  that can be expressed as a linear combination of the columns of  $\mathbf{V}$ 

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• Apply Pythagorus to  $\vec{y} - \vec{z}$  for arbitrary  $\vec{z} \in \mathbb{R}^n$ :

$$\begin{aligned} \|\vec{y} - \vec{z}\|_{2}^{2} &= \|\mathbf{V}\mathbf{V}^{T}(\vec{y} - \vec{z})\|_{2}^{2} + \|\vec{y} - \vec{z} - \mathbf{V}\mathbf{V}^{T}(\vec{y} - \vec{z})\|_{2}^{2} \\ &= \|\mathbf{V}\mathbf{V}^{T}\vec{y} - \mathbf{V}\mathbf{V}^{T}\vec{z}\|_{2}^{2} + \|\vec{y} - \mathbf{V}\mathbf{V}^{T}\vec{y} + \vec{z} - \mathbf{V}\mathbf{V}^{T}\vec{z}\|_{2}^{2}. \end{aligned}$$

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• If  $\vec{z} = \mathbf{V}\vec{c}$  for some  $\vec{c} \in \mathbb{R}^k$ , then  $\mathbf{V}\mathbf{V}^T\vec{z} = \mathbf{V}\mathbf{V}^T\mathbf{V}\vec{c} = \mathbf{V}\vec{c} = \vec{z}$  and the above simplifies to

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• To minimize this, set  $\vec{z} = \mathbf{V}\mathbf{V}^T\vec{y}$ .

If  $\vec{x_1}, \ldots, \vec{x_n}$  are close to a *k*-dimensional subspace  $\mathcal{V}$  with orthonormal basis  $\mathbf{V} \in \mathbb{R}^{d \times k}$ , the data matrix can be approximated as  $\mathbf{XVV}^T$  and  $\mathbf{XV}$  gives optimal embedding of  $\mathbf{X}$  in  $\mathcal{V}$ . How do we find  $\mathcal{V}$  (equivalently  $\mathbf{V}$ )?

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$$\|\mathbf{X} - \mathbf{X}\mathbf{V}\mathbf{V}^{T}\|_{F}^{2} = \|\mathbf{X}^{T} - \mathbf{V}\mathbf{V}^{T}\mathbf{X}^{T}\|_{F}^{2}$$
  
$$= \sum_{i=1}^{n} \|\vec{x}_{i} - \mathbf{V}\mathbf{V}^{T}\vec{x}_{i}\|_{2}^{2}$$
  
$$= \sum_{i=1}^{n} \|\vec{x}_{i}\|_{2}^{2} - \|\mathbf{V}\mathbf{V}^{T}\vec{x}_{i}\|_{2}^{2}$$

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So minimizing  $\|\mathbf{X} - \mathbf{X}\mathbf{V}\mathbf{V}^T\|_F^2$  is the same as maximizing

$$\sum_{i} \| \mathbf{V} \mathbf{V}^{\mathsf{T}} \vec{x}_{i} \|_{2}^{2} = \sum_{i} \vec{x}_{i}^{\mathsf{T}} \mathbf{V} \mathbf{V}^{\mathsf{T}} \mathbf{V} \mathbf{V}^{\mathsf{T}} \vec{x}_{i} = \sum_{i} \| \mathbf{V}^{\mathsf{T}} \vec{x}_{i} \|_{2}^{2}$$

 $\boldsymbol{V}$  minimizing  $\|\boldsymbol{X}-\boldsymbol{X}\boldsymbol{V}\boldsymbol{V}^{\mathcal{T}}\|_{\textit{F}}^2$  is given by:

$$\underset{\text{orthonormal } \mathbf{V} \in \mathbb{R}^{d \times k}}{\arg \max} \sum_{i=1}^{n} \|\mathbf{V}^{\mathsf{T}} \vec{x_i}\|_2^2 = \sum_{j=1}^{k} \sum_{i=1}^{n} \langle \vec{v_j}, \vec{x_i} \rangle^2$$

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Surprisingly, can find the columns of **V**,  $\vec{v_1}, \ldots, \vec{v_k}$  greedily.

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$$\vec{v}_{k} = \arg \max_{\vec{v} \text{ with } \|v\|_{2}=1, \ \langle \vec{v}, \vec{v}_{j} \rangle = 0 \ \forall j < k} \vec{v}^{T} \mathbf{X}^{T} \mathbf{X} \vec{v}$$

**V** minimizing  $\|\mathbf{X} - \mathbf{X}\mathbf{V}\mathbf{V}^T\|_F^2$  is given by:

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. . .

These are exactly the top k eigenvectors of  $\mathbf{X}^T \mathbf{X}$ .

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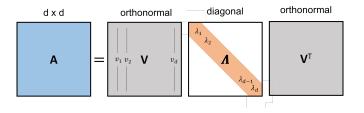
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Yields eigendecomposition:  $AVV^T = A = V\Lambda V^T$  where the first inequality follows since rows of **A** are in span of the eigenvectors.



Typically order the eigenvectors in decreasing order:

$$\lambda_1 \geq \lambda_2 \geq \ldots \geq \lambda_d$$

**Courant-Fischer Principal:** For symmetric **A**, the eigenvectors are given via the greedy optimization:

$$\vec{v}_1 = \underset{\vec{v} \text{ with } \|v\|_2=1}{\arg \max} \vec{v}^T \mathbf{A} \vec{v}.$$
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$$\cdots$$
$$\vec{v}_d = \arg \max \qquad \vec{v}^T \mathbf{A} \vec{v}.$$

$$\vec{v}$$
 with  $\|v\|_2 = 1$ ,  $\langle \vec{v}, \vec{v}_j \rangle = 0 \quad \forall j < d$ 

• 
$$\vec{v}_j^T \mathbf{A} \vec{v}_j = \lambda_j \cdot \vec{v}_j^T \vec{v}_j = \lambda_j$$
, the  $j^{th}$  largest eigenvalue.

**Courant-Fischer Principal:** For symmetric **A**, the eigenvectors are given via the greedy optimization:

$$\vec{v}_{1} = \underset{\vec{v} \text{ with } \|v\|_{2}=1}{\arg \max} \vec{v}^{T} \mathbf{A} \vec{v}.$$
$$\vec{v}_{2} = \underset{\vec{v} \text{ with } \|v\|_{2}=1, \ \langle \vec{v}, \vec{v}_{1} \rangle = 0}{\arg \max} \vec{v}^{T} \mathbf{A} \vec{v}.$$
$$\cdots$$
$$\vec{v}_{d} = \underset{\vec{v} \text{ with } \|v\|_{2}=1, \ \langle \vec{v}, \vec{v}_{i} \rangle = 0}{\arg \max} \vec{v}^{T} \mathbf{A} \vec{v}.$$

•  $\vec{v}_j^T \mathbf{A} \vec{v}_j = \lambda_j \cdot \vec{v}_j^T \vec{v}_j = \lambda_j$ , the  $j^{th}$  largest eigenvalue.

The first k eigenvectors of X<sup>T</sup>X (corresponding to the largest k eigenvalues) are exactly the directions of greatest "variance" in X that we use for low-rank approximation. We'll talk more about this next time.