# COMPSCI 514: ALGORITHMS FOR DATA SCIENCE

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Lecture 19

#### **Spectral Graph Partitioning**

- Focus on separating graphs with small but relatively balanced cuts.
- Connection to second smallest eigenvector of graph Laplacian.
- Today: Provable guarantees for stochastic block model.

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- Haven't given formal guarantees; it's difficult for general input graphs. But can consider randoms "natural" graphs...
- **Common Approach:** Give a natural generative model for random inputs and analyze how the algorithm performs on inputs drawn from this model. Can be used to justify  $\ell_2$  linear regression, *k*-means clustering, etc.

# STOCHASTIC BLOCK MODEL

**Stochastic Block Model (Planted Partition Model):** Let  $G_n(p,q)$  be a distribution over graphs on *n* nodes, split randomly into two groups *B* and *C*, each with n/2 nodes.

- Any two nodes in the same group are connected with probability *p* (including self-loops).
- Any two nodes in different groups are connected with prob. *q* < *p*.
- Connections are independent.



### LINEAR ALGEBRAIC VIEW

Let G be a stochastic block model graph drawn from  $G_n(p, q)$ .

• Let  $\mathbf{A} \in \mathbb{R}^{n \times n}$  be the adjacency matrix of G, ordered in terms of group ID.



 $G_n(p, q)$ : stochastic block model distribution. *B*, *C*: groups with n/2 nodes each. Connections are independent with probability *p* between nodes in the same group, and probability *q* between nodes not in the same group.

Letting G be a stochastic block model graph drawn from  $G_n(p,q)$  and  $\mathbf{A} \in \mathbb{R}^{n \times n}$  be its adjacency matrix.  $(\mathbb{E}[\mathbf{A}])_{i,j} = p$  for i, j in same group,  $(\mathbb{E}[\mathbf{A}])_{i,j} = q$  otherwise.



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What is rank( $\mathbb{E}[\mathbf{A}]$ )? What are the eigenvectors and eigenvalues of  $\mathbb{E}[\mathbf{A}]$ ?

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- Second eigenvector of A is close to [1, 1, 1, ..., -1, -1, -1] and gives a good estimate of the communities.



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When rows/columns aren't sorted by ID, second eigenvector is e.g.,  $[1, -1, 1, -1, \ldots, 1, 1, -1]$  and entries give community ids.

Letting G be a stochastic block model graph drawn from  $G_n(p, q)$ ,  $\mathbf{A} \in \mathbb{R}^{n \times n}$  be its adjacency matrix and  $\mathbf{L}$  be its Laplacian, what are the eigenvectors and eigenvalues of  $\mathbb{E}[\mathbf{L}]$ ? Letting G be a stochastic block model graph drawn from  $G_n(p, q)$ ,  $\mathbf{A} \in \mathbb{R}^{n \times n}$  be its adjacency matrix and  $\mathbf{L}$  be its Laplacian, what are the eigenvectors and eigenvalues of  $\mathbb{E}[\mathbf{L}]$ ?

$$\mathbb{E}[\mathbf{L}] = \mathbb{E}[\mathbf{D}] - \mathbb{E}[\mathbf{A}] = \left(\frac{n(p+q)}{2}\right)\mathbf{I} - \mathbb{E}[\mathbf{A}]$$

and so if  $\mathbb{E}[\mathbf{A}]\vec{x} = \lambda \vec{x}$  then

$$\mathbb{E}[\mathbf{L}]\vec{x} = (n(p+q)/2 - \lambda)\vec{x}$$

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Therefore the first and second eigenvalues of  $\mathbb{E}[A]$  are the second and first eigenvectors of  $\mathbb{E}[L].$ 

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How do we show that a matrix is close to its expectation? Matrix concentration inequalities.

- Analogous to scalar concentration inequalities like Markovs, Chebyshevs, Bernsteins.
- Random matrix theory is a very recent and cutting edge subfield of mathematics that is being actively applied in computer science, statistics, and ML.

**Matrix Concentration Inequality:** If  $p \ge O\left(\frac{\log^4 n}{n}\right)$ , then with high probability

$$\|\mathbf{A} - \mathbb{E}[\mathbf{A}]\|_2 \leq O(\sqrt{pn}).$$

where  $\|\cdot\|_2$  is the matrix spectral norm (operator norm).

For any  $\mathbf{X} \in \mathbb{R}^{n \times d}$ ,  $\|\mathbf{X}\|_2 = \max_{z \in \mathbb{R}^d: \|z\|_2 = 1} \|\mathbf{X}z\|_2$ .

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For the stochastic block model application, we want to show that the second eigenvectors of **A** and  $\mathbb{E}[\mathbf{A}]$  are close. How does this relate to their difference in spectral norm?

**Davis-Kahan Eigenvector Perturbation Theorem:** Suppose  $\mathbf{A}, \overline{\mathbf{A}} \in \mathbb{R}^{d \times d}$  are symmetric with  $\|\mathbf{A} - \overline{\mathbf{A}}\|_2 \leq \epsilon$  and eigenvectors  $v_1, v_2, \ldots, v_d$  and  $\overline{v}_1, \overline{v}_2, \ldots, \overline{v}_d$ . Letting  $\theta(v_i, \overline{v}_i)$  denote the angle between  $v_i$  and  $\overline{v}_i$ , for all i:

$$\sin[\theta(v_i, \bar{v}_i)] \le \frac{\epsilon}{\min_{j \neq i} |\lambda_i - \lambda_j|}$$

where  $\lambda_1, \ldots, \lambda_d$  are the eigenvalues of  $\overline{\mathbf{A}}$ .

The errors get large if there's eigenvalues with similar magnitudes.

Claim 1 (Matrix Concentration): For  $p \ge O\left(\frac{\log^4 n}{n}\right)$ ,  $\|\mathbf{A} - \mathbb{E}[\mathbf{A}]\|_2 \le O(\sqrt{pn}).$ 

Claim 2 (Davis-Kahan): For  $p \ge O\left(\frac{\log^4 n}{n}\right)$ ,

$$\sin\theta(v_2,\bar{v}_2) \leq \frac{O(\sqrt{\rho n})}{\min_{j\neq 2}|\lambda_2 - \lambda_j|}$$

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**Recall:**  $\mathbb{E}[\mathbf{A}]$  has eigenvalues  $\lambda_1 = \frac{(p+q)n}{2}$ ,  $\lambda_2 = \frac{(p-q)n}{2}$ ,  $\lambda_i = 0$  for  $i \ge 3$ .

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- So they differ in sign in at most  $O\left(\frac{p}{(p-q)^2}\right)$  positions.

**Upshot:** If *G* is a stochastic block model graph with adjacency matrix **A**, if we compute its second large eigenvector  $v_2$  and assign nodes to communities according to the sign pattern of this vector, we will correctly assign all but  $O\left(\frac{p}{(p-q)^2}\right)$  nodes.

